



SAT Encodings for Sudoku

Bug Catching in 2006 Fall

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Agenda

- **Introduction**
- Background and Previous Encodings
- Optimized Encoding
- Experimental Results
- Conclusions

What is Sudoku ?

Problem

		6	1		2	5		
	3	9				1	4	
				4				
9		2		3		4		1
	8						7	
1		3		6		8		9
				1				
	5	4				9	1	
		7	5		3	2		

Given a problem, the objective is to find a **satisfying assignment** w.r.t. Sudoku rules.



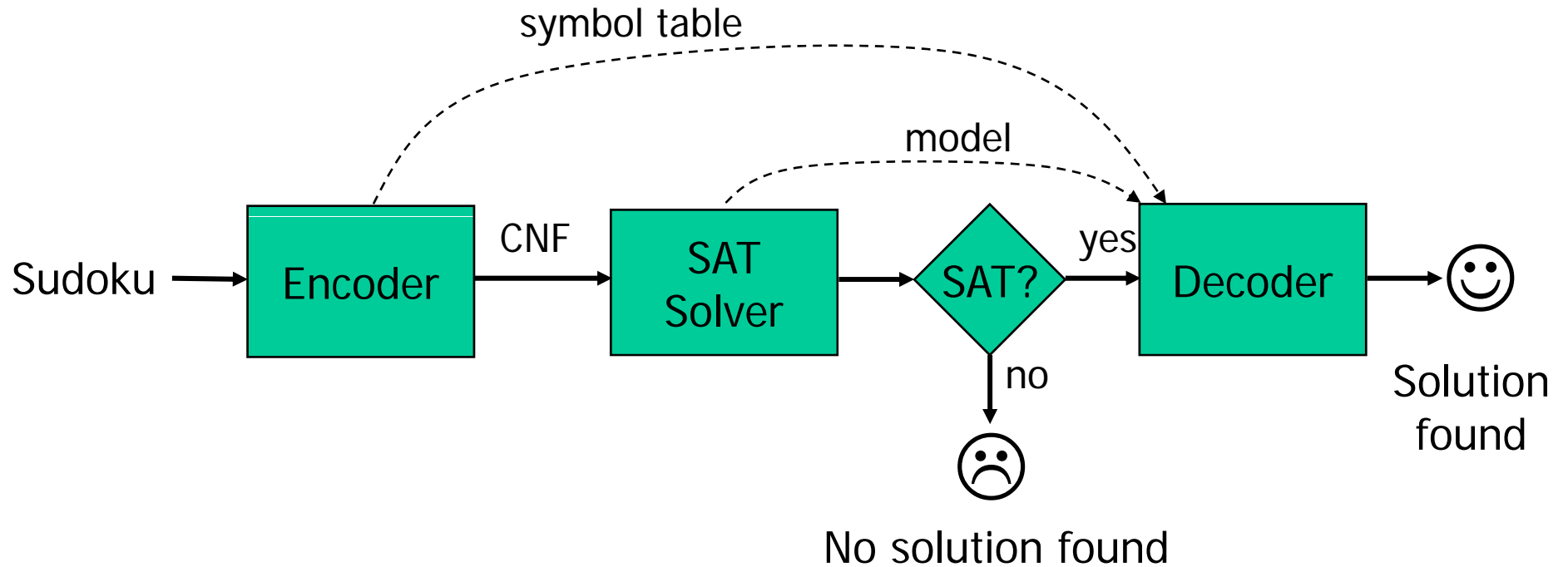
Solution

8	4	6	1	7	2	5	9	3
7	3	9	6	5	8	1	4	2
5	2	1	3	4	9	7	6	8
9	6	2	8	3	7	4	5	1
4	8	5	9	2	1	3	7	6
1	7	3	4	6	5	8	2	9
2	9	8	7	1	4	6	3	5
3	5	4	2	8	6	9	1	7
6	1	7	5	9	3	2	8	4

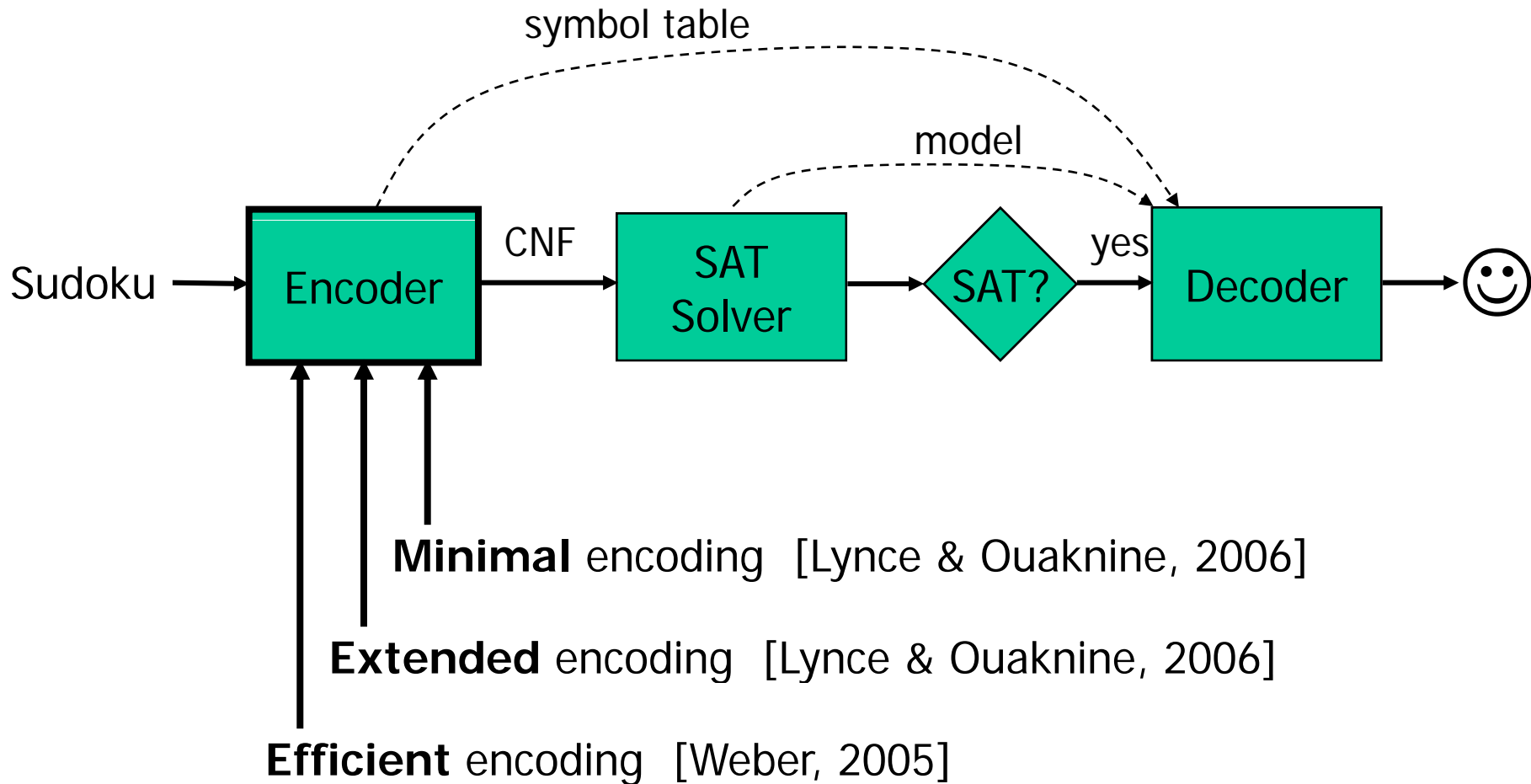
Sudoku rules

- ✓ There is a number in each **cell**.
- ✓ A number appears once in each **row**.
- ✓ A number appears once in each **column**.
- ✓ A number appears once in each **block**.

Sudoku as SAT Problem



Previous Encodings

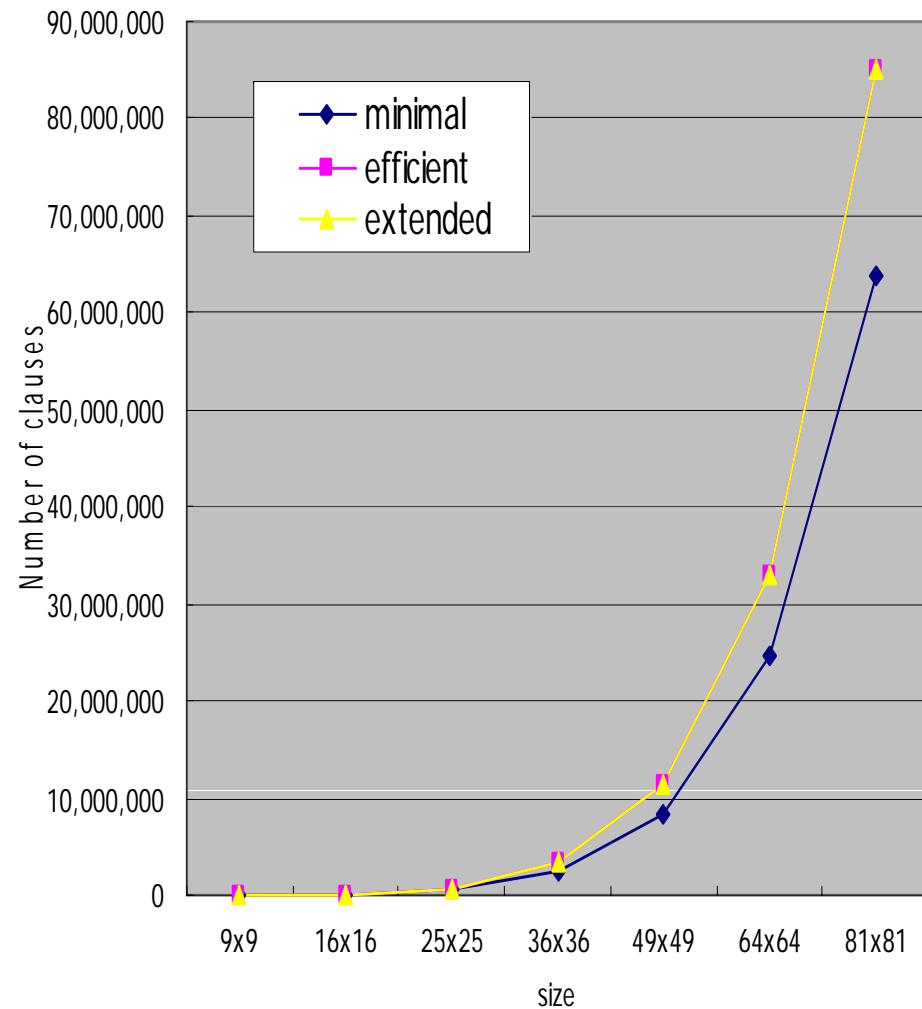


Analysis of Previous Encodings

Encoding	Number of Variables	Number of Clauses
Minimal	N^3	$N * N + \left(N * N * \left(\frac{N * (N - 1)}{2} \right) \right) * 3 + k$
Efficient	N^3	$N * N + \left(N * N * \left(\frac{N * (N - 1)}{2} \right) \right) * 4 + k$
Extended	N^3	$\left(N * N + N * N * \left(\frac{N * (N - 1)}{2} \right) \right) * 4 + k$

Exponential Growth in Clauses



size	minimal	efficient	extended
9x9	8829	11745	11988
16x16	92416	123136	123904
25x25	563125	750625	752500
36x36	2450736	3267216	3271104
49x49	8473129	11296705	11303908
64x64	24776704	33034240	33046528
81x81	63779481	85037121	85056804



Experimental Results

		minimal encoding			efficient encoding			extended encoding		
size	level	vars	clauses	time	vars	clauses	time	vars	clauses	time
9x9	easy	729	8854	0.00	729	11770	0.00	729	12013	0.00
9x9	hard	729	8859	0.00	729	11775	0.00	729	12018	0.00
16x16	easy	4096	92520	0.10	4096	123240	0.09	4096	124008	0.01
16x16	hard	4096	92514	0.46	4096	123234	0.21	4096	124002	0.01
25x25	easy	15625	563417	9.07	15625	750917	17.48	15625	752792	0.07
25x25	hard	15625	563403	time	15625	750903	time	15625	752778	0.21
36x36	easy	46656	2451380	time	46656	3267860	time	46656	3271748	0.50
36x36	hard	46656	2451400	time	46656	3267880	time	46656	3271768	0.67
49x49	easy	117649	8474410	time	117649	11297986	time	117649	11305189	1.47
64x64	easy	262144	24779088	stack	262144	33036624	stack	262144	33048912	stack
81x81	easy	531441	63783464	stack	531441	85041104	stack	531441	85060787	stack

Experimental Results

		minimal encoding			efficient encoding			extended encoding		
size	level	vars	clauses	time	vars	clauses	time	vars	clauses	time
9x9	easy	729	8854	0.00	729	11770		29	12013	0.00
9x9	hard	729	8859	0.00	729	11775		29	12018	0.00
16x16	easy	4096	92520	0.10	4096	123240		96	124008	0.01
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64x64	easy	262144			262144	33048912	stack	262144	33048912	stack
81x81	easy	531441			531441	85060787	stack	531441	85060787	stack

Solution found

No solution found

Motivations

- Sudoku was regarded as SAT problem
 - W Weber, [A SAT-based Sudoku Solver](#), Nov. 2005.
 - Lynce & Ouaknine, [Sudoku as a SAT Problem](#), Jan. 2006.
 - ➔ Extended encoding shows the best performance in our experiments
- Problems in previous works
 - Too many clauses are generated (e.g. 85,056,804 clauses)
 - Thus, large size puzzles are not solved
 - ➔ The extended encoding must be **optimized** for large size puzzles
- Our objectives
 - To propose an **optimization** for the extended encoding

Agenda

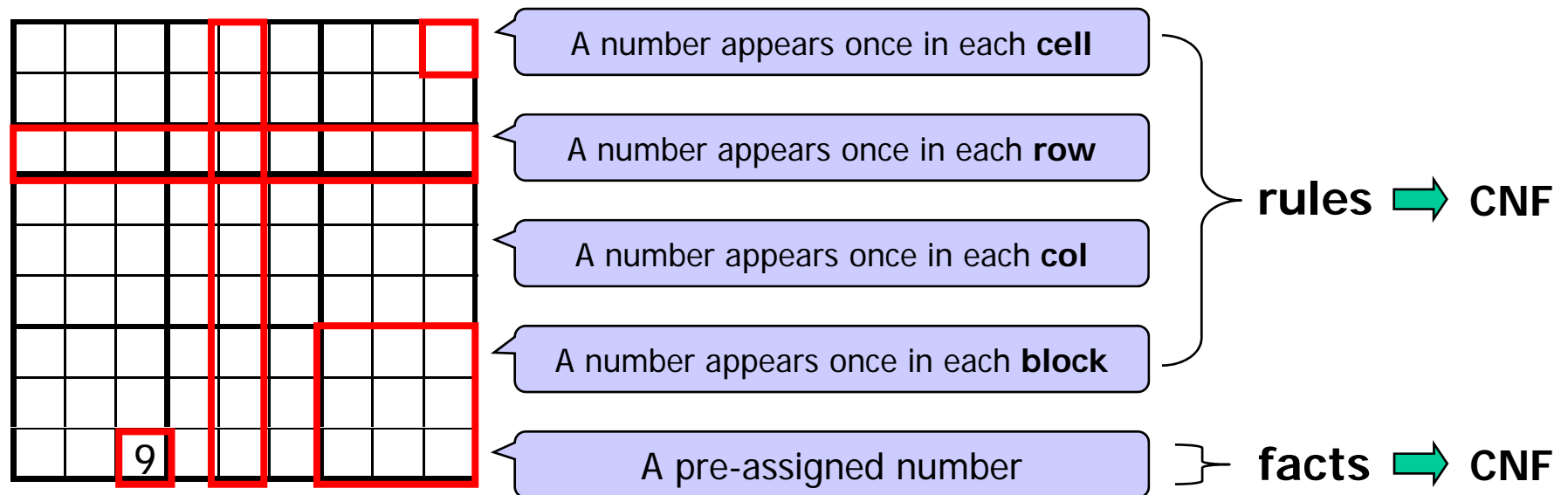
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Encoding

- Knowledge compilation into a **target language**

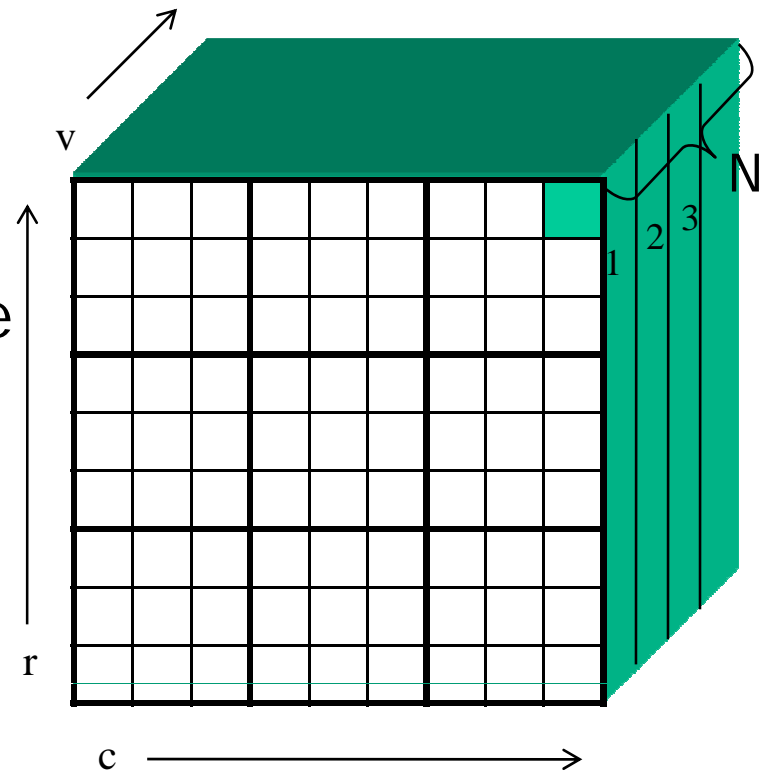
problem knowledge → CNF

- Knowledge about Sudoku

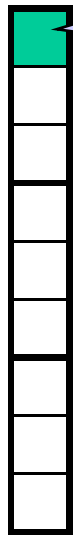


Variables

- Each cell has one number from 1..N
 - $[1,1]=1$ or $[1,1]=2$ or or $[1,1]=N$
 - Each cell needs N boolean variables to consider all cases
- Total number of variables
 - N^3
- Boolean variable name as a triple
 - (r,c,v) (i.e., x_{rcv}) iff $[r,c] = v$
 - $\neg(r,c,v)$ (i.e., $\neg x_{rcv}$) iff $[r,c] \neq v$



Cell Rule → CNF



A number appears once in each **cell**

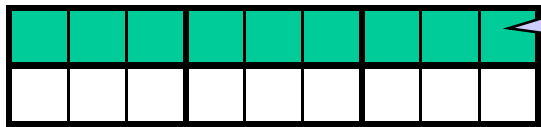
There is **at least** one number in each cell (definedness)

$$Cell_d = \bigwedge_{r=1}^N \bigwedge_{c=1}^N \bigvee_{v=1}^N (r, c, v)$$

There is **at most** one number in each cell (uniqueness)

$$Cell_u = \bigwedge_{r=1}^N \bigwedge_{c=1}^N \bigwedge_{v_i=1}^{N-1} \bigwedge_{v_j=v_i+1}^N \neg((r, c, v_i) \wedge (r, c, v_j))$$

Row Rule → CNF



A number appears once in each **row**

Each number appears **at least** once in each row (definedness)

$$Row_d = \bigwedge_{r=1}^N \bigwedge_{v=1}^N \bigvee_{c=1}^N (r, c, v)$$

Each number appears **at most** once in each row (uniqueness)

$$Row_u = \bigwedge_{r=1}^N \bigwedge_{v=1}^N \bigwedge_{c_i=1}^{N-1} \bigwedge_{c_j=c_i+1}^N \neg((r, c_i, v) \wedge (r, c_j, v))$$

Column Rule → CNF

A number appears once in each **column**

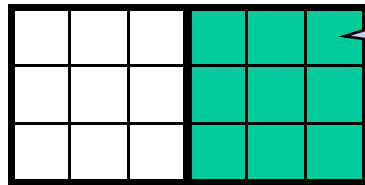
Each number appears **at least** in each column (**definedness**)

$$Col_d = \bigwedge_{c=1}^N \bigwedge_{v=1}^N \bigvee_{r=1}^N (r, c, v)$$

Each number appears **at most** in each column (**uniqueness**)

$$Col_u = \bigwedge_{c=1}^N \bigwedge_{v=1}^N \bigwedge_{r_i=1}^{N-1} \bigwedge_{r_j=r_i+1}^N \neg((r_i, c, v) \wedge (r_j, c, v))$$

Block Rule → CNF



A number appears once in each **block**

Each number appears **at least** once in each block (definedness)

$$Block_d = \bigwedge_{r_{offs}=1}^{subN} \bigwedge_{c_{offs}=1}^{subN} \bigwedge_{v=1}^N \bigvee_{r=1}^{subN} \bigvee_{c=1}^{subN} (r_{offs} * subN + r, c_{offs} * subN + c, v)$$

Each number appears **at most** once in each block (uniqueness)

$$Block_u = \bigwedge_{r_{offs}=1}^{subN} \bigwedge_{c_{offs}=1}^{subN} \bigwedge_{v=1}^N \bigwedge_{r=1}^N \bigwedge_{c=r+1}^N \\ \neg((r_{offs} * subN + (r \bmod subN), c_{offs} * subN + (r \bmod subN), v) \\ \wedge (r_{offs} * subN + (c \bmod subN), c_{offs} * subN + (c \bmod subN), v))$$

Pre-Assigned Fact \rightarrow CNF

		3

A pre-assigned number

As a constant; the number is never changed

It can be represented as a **unit clause**

$$Assigned = \bigwedge_{i=1}^k \{ (r, c, a) \mid \exists_{1 \leq a \leq N} \bullet [r, c] = a \}$$

where k is a number of pre - assigned numbers

Previous Encodings

Minimal encoding [Lynce & Ouaknine, 2006]

$$\phi = Cell_d \cup Row_u \cup Col_u \cup Block_u \cup Assigned$$

sufficient to characterize the puzzle

Extended encoding [Lynce & Ouaknine, 2006]

$$\phi = Cell_d \cup Cell_u \cup Row_d \cup Row_u \cup Col_d \cup Col_u \\ \cup Block_d \cup Block_u \cup Assigned$$

minimal encoding with redundant clauses

Efficient encoding [Weber, 2005]

$$\phi = Cell_d \cup Cell_u \cup Row_u \cup Col_u \cup Block_u \cup Assigned$$

between minimal encoding and extended encoding

Analysis (Recap)

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Minimal	N^3	$N * N + \left(N * N * \left(\frac{N * (N - 1)}{2} \right) \right) * 3 + k$
Efficient	N^3	$N * N + \left(N * N * \left(\frac{N * (N - 1)}{2} \right) \right) * 4 + k$
Extended	N^3	$\left(N * N + N * N * \left(\frac{N * (N - 1)}{2} \right) \right) * 4 + k$