The Spin Model Checker : Part I

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Hierarchy of SW Coverage Criteria



Model Checker Analyzes All Possible Scheduling

```
active type A() {
                                                             x:0
  byte x;
  again:
                                                             X:1
      X++;
      goto again;
                                                             x:2
  }
                                                            x:255
   active type A() {
   byte x;
   again:
                                                                        (x:0,y:
                                                           х:0,у
                                                                                         x:0,y:255
       X++;
      goto again;
                                                                        (x:1,y:)
                                                          (x:1,y:
                                                                                         (;1,y:255
                                                                                         :2,y:255
                                                          (x:2,y:
                                                                        (x:2,y:
   active type B() {
   byte y;
   again:
                                                                                        <del>*</del>:255,y:255
       y++;
                                                           (:255,y)0
     goto again;
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```

Overview of the Spin Architecture



- A few characteristics of Spin
 - Promela allows a finite state model only
 - Asynchronous execution
 - Interleaving semantics for concurrency
 - 4 2-way process communication
 - Non-determinism

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Promela provides (comparatively) rich set of constructs such as variables and message passing, dynamic creation of processes, etc

Tcl GUI of SPIN (ispin.tcl): Edit Window

•	mobile1.pml							
	Spin Version 6.2.7 2 March 2014 :: iSpin Version 1.1.0 7 June 2	2012						
Edit/Viev	ew Simulate / Replay Verification Swarm Run <help> Save Session Resto</help>	ore Session <quit></quit>						
Open	ReOpen Save Save As Syntax Check Redundancy Check Symbol Table Find:							
1 2 3 4 5 6 2). 7 8 9 10 11 12 13 14 15 16 17 blue }; 18 19 20	<pre>/* * Model of cell-phone handoff strategy in a mobile network. * A translation from the pi-calculus description of this * model presented in: * Fredrik Orava and Joachim Parrow, 'An algebraic verification * of a mobile network,' Formal aspects of computing, 4:497-543 (199 * For more information on this model, email: joachim@it.kth.se * See also the simplified version of this model in mobile2 * * to perform the verification: * \$ spin -a mobile1 * \$ cc -o pan pan.c * \$ pan -a */ mtype = { data, ho_crmd, ho_com, ho_acc, ho_fail, ch_rel, white, red, byte a_id, p_id; /* ids of processes refered to in the property */ </pre>	mata View zoom in zoom out Select: p_CC p_HC o_MSC p_BS p_MS p_P p_Q System p_bot aim_ltl_0						
3 gcc - 4 ./pan 5 C:/sp	-o pan pan.c n -D dot > dot.tmp pin/mobile1.pml:1	^						
6 spin -o3 -a mobile1.pml Itl Itl_0: (! ([] (<> (((BS[a_id]@progress)) ((BS[p_id]@progress))))))) ([] ((! (<> (inp?[red]))) (<> (out?[red])))) 7 gcc -o pan pan.c 8 ./pan -D dot > dot.tmp								

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Tcl GUI of SPIN (ispin.tcl): Verification Window

9		2th	eads.pml		- 🗆 ×					
Spin Version 6.2.7 2 March 2014 :: iSpin Version 1.1.0 7 June 2012										
Edit/View Simulate / Replay Verification	Swarm Run <help></help>	Save Session Restore S	ession <quit></quit>							
Safety Storage Mode		Search Mode	Remove	B	Remove					
	e	 depth-first search 	Advanced: Error Trappin	ng Advance	Advanced: Parameters					
🔽 + invalid endstates (deadlock) 🛛 🗖 + mini	mized automata (slow)	✓ + partial order reduction	n on't stop at errors	Physical Memory Available (ii	in Mbytes): 1024 explain					
✓ + assertion violations	pse compression	+ bounded context sw	tching O stop at error nr: 1	Estimated State Space Size (s	states x 10.42): 1000 explain					
□ + xr/xs assertions	npact C bitstate/supertrace	with bound: 0	✓ save all error-trails	Estimated State Space Size (s	states x 10 ⁻¹ 5). 1000 explain					
Liveness	Never Claims	+ iterative search for s	nort trail 🔲 add complexity profiling	Maximum Search Depth (step	ps): 10000 explain					
C non-progress cycles	se a never claim or Itl property	O breadth-first search	compute variable ranges	Nr of hash-functions in Bitsta	ate mode: 3 explain					
C acceptance cycles C use claim	1	🔽 + partial order reduction	n A Full Channel	Size for Minimized Automato	on 100 explain					
enforce weak fairness constraint claim name	(opt):	report unreachable code	blocks new msgs	Extra Verifier Generation Opti	ions: explain					
		Save Result in: pan.out	C loses new msgs	Extra Compile-Time Directive	explain					
Run	Stop		State Tables Clear	Help Extra Run-Time Options:	explain					
Run Stop 1 byte x=0, y=0, z =0; 2 active proctype p() { 4 $x=y+1;$ 5 $y=z+1;$ 6 $z=x+1;$ 7 } 8 9 9 active proctype q() { 10 $y=z+1;$ 11 $z=x+1;$ 12 $x=y+1;$ 13 } 14 15 15 init { 16 timeout-> printf("x:%d,y:%d,z:%d\n",x,y,z); 17 assert(0); 18 } 19 20 20 /* 21 spin -a 2threads.pml 22 gcc -DMEMLIM=1024 -O2 -DXUSAFE -DSAFETY -DNOCLAIM -		AlM -w -o pa	 Hull statespace search tor: never claim - (not selected) assertion violations + cycle checks - (disabled by -DSAFETY) invalid end states + State-vector 24 byte, depth reached 12, errors: 11 115 states, stored 7 states, matched 122 transitions (= stored+matched) 0 atomic steps hash conflicts: 0 (resolved) Stats on memory usage (in Megabytes): 0.004 equivalent memory usage for states (stored*(State-vector + overhead))) 0.292 actual memory usage for states 64.000 memory used for hash table (-w24) 0.343 memory usage for DFS stack (-m10000) 64.539 total actual memory usage 							



Tcl GUI of SPIN (ispin.tcl): Simulation Window

				2thre	ads.pml					_ □	א נ
		Sp	in Version 6.2	.7 2 March 201	4 :: iSpin Version 1	1.1.0 7	7 June 2	2012			
Edit/View	Simulate / Replay	Verification	Swarm Run	<help> Sav</help>	e Session Re	store Se	ession	<quit></quit>			
	Mode		A	Full Channel	Output Filterin	g (reg. e	xps.)	(Re)Run	Background co	mmand execu	uted:
C Random	n, with seed:	123	💿 bloc	cks new messages	process ids:			(spin -p -s -r -X	-v -n123 -l -	g -k C
O Interacti	O Interactive (for resolution of all nondeterminism)			C loses new messages queue ids:				Stop	:/Dropbox/classes/Spring14-cs 92B/2threads/2threads.pml5.tr		
Ouided,	with trail: C:/Dropbox/	classes/Spring brow	se 🔽 MSC+	⊦stmnt	var names:			Rewind	-u10000 2threa	ids.pml	
initial step	os skipped:	0	MSC max	c text width 20	tracked variable			Step Forward			
maximum	number of steps:	10000	MSC upo	date delay 25		• 		Chan Dealaward	-		
Track D	ata Values (this can be s	low)			track scaling:		_	Step васкward		Save in: m	nsc.ps
1	int x=0, y=0, z=0;					ÎT	I MSC -	<u> </u>	q:1:1		í
3	active proctype p() {	•					2	р х =	(v+1) y = (z+	1)	
4	x=y+1; v=z+1						3		z = (x+	·1)	
6	z=x+1;						4	у =	(z+1)		
7	}						5		(y+	1)	1.2
9	active proctype q() {						7	2 -	(X+1)	time.	.⊥.∠ eout)
10	y=z+1;						8			printf('x:%d,	y:%d
12	x=y+1;										
13	1					¥]	<				>
[variable tep 1] y = 1	e values, s ⊥ 1: 2: 3: 4: 5: 6: 7: x:5,y: 8: spin: spin: #proo 9:	proc 1 (q:1 proc 0 (p:1 proc 1 (q:1 proc 0 (p:1 proc 1 (q:1 proc 1 (q:1 proc 2 (:ini ;4,z:6 proc 2 (:ini 2threads.pml:1 text of failed ass cesses: 3 proc 2 (:ini) 2threads.) 2threads.) 2threads.) 2threads.) 2threads.) 2threads.) 2threads. t::1) 2thread t::1) 2thread f, Error: ass sertion: ass t::1) 2thread	pml:10 (state 1) pml:4 (state 1) pml:11 (state 2) pml:5 (state 2) pml:12 (state 3) pml:6 (state 3) ds.pml:16 (state ds.pml:16 (state sertion violated ert(0) ds.pml:17 (state	[y = (z+1)] x = (y+1)] [z = (x+1)] y = (z+1)] [x = (y+1)] z = (x+1)] = 1) [(timeout)] = 2) [printf('x:%d	,y:%d,z	z:%d\\r	n',x,y,z)]	^ [que	ues, step	1]

Overview of the Promela



- Similar to C syntax but simplified
 - No pointer
 - No real datatype such as float or real
 - 4 No functions
- Processes are communicating with each other using
 - Global variables
 - 4 Message channels
- Process can be dynamically created
- Scheduler executes one process at a time using interleaving semantics

active[2] proctype A() {
 byte x;
 printf("A%d is starting\n");
}

proctype B() {
 printf("B is starting\n");
}

```
Init {
run B();
}
```

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run() operator creates a process and returns a newly created process ID

There are 6 possible outcomes due to non-deterministic scheduling
 4 A0.A1.B, A0.B.A1
 4 A1.A0.B, A1.B.A0
 4 B.A0.A1, B.A1.A0
 In other words, process

In other words, process creation may not immediately start process execution

Basic types

- \rm bit
- 🜲 bool
- Byte (8 bit unsigned integer)
- short (16 bits signed integer)
- Int (32 bits signed integer)
- Arrays
 - bool x[10];
- Records

- typedef R { bit x; byte y;}
- Default initial value of variables is 0
- Most arithmetic (e.g.,+,-), relational (e.g. >,==) and logical operators of C are supported
 - bitshift operators are supported too.

- Promela spec generates only a finite state model because
 - Max # of active process <= 255</p>
 - Each process has only finite length of codes
 - Each variable is of finite datatype
 - All message channels have bounded capability <= 255</p>



- Each Promela statement is either
 - **4** executable:
 - Blocked

- There are six types of statement
 - Assignment: always executable
 - Ex. x=3+x, x=run A()
 - Print: always executable
 - Ex.printf("Process %d is created.\n",_pid);
 - Assertion: always executable
 - Ex. assert(x + y == z)
 - Expression: depends on its value
 - Ex. x+3>0, 0, 1, 2
 - Ex.skip, true
 - Send: depends on buffer status
 - Ex. ch1!m is executable only if ch1 is not full
 - Receive: depends on buffer status
 - Ex. ch1?m is executable only if ch1 is not empty

An expression is also a statement It is executable if it evaluates to non-zero 41 : always executable 41<2:always executable</p> 4x<0: executable only when x < 0 +x-1:executable only when x !=0If an expression statement in blocked, it remains blocked until other process changes the condition an expression e is equivalent to while(!e); in C



assert(expr)

- assert is always executable
- #If expr is 0, SPIN detects this violation
- assert is most frequently used checking method, especially as a form of invariance
 - ex. active proctype inv() { assert(x== 0);}
 Note that inv() is equivalent to [] (x==0) in LTL with thanks to interleaving semantics



Generation of all possible interleaving scenarios



Therefore, just a single assert(x==0) statement in Inv() can check if x==0 all the time



Program Execution Control

- Promela provides low-level control mechanism, i.e., goto and label as well as if and do
- Note that non-deterministic selection is supported
- else is predefined variable which becomes true if all guards are false; false otherwise



Critical Section Example

```
[root@moonzoo spin_test]# ls
                                           crit.pml
                                            [root@moonzoo spin_test]# spin -a crit.pml
                                            [root@moonzoo spin_test]# ls
                                            crit.pml pan.b pan.c pan.h pan.m pan.t
                                            [root@moonzoo spin_test]# gcc pan.c
bool lock;
                                            [root@moonzoo spin_test]# a.out
byte cnt;
                                            pan: assertion violated (cnt<=1) (at depth 8)
                                            pan: wrote crit.pml.trail
active[2] proctype P() {
                                           Full statespace search for:
     !lock -> lock=true;
                                                                    - (none specified)
                                                never claim
                                                assertion violations
     cnt=cnt+1:
                                                                      +
                                                acceptance cycles - (not selected)
     printf("%d is in the crt sec!\n",_pid);
                                                invalid end states
                                                                     +
     cnt=cnt-1;
                                           State-vector 36 byte, depth reached 16, errors: 1
     lock=false;
                                               119 states, stored
                                               47 states, matched
                                               166 transitions (= stored+matched)
active proctype Invariant() {
                                                0 atomic steps
     assert(cnt <= 1);
                                           hash conflicts: 0 (resolved)
                                           4.879 memory usage (Mbyte)
}
                                            [root@moonzoo spin_test]# ls
                                           a.out crit.pml crit.pml.trail pan.b pan.c pan.h
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                                           pan.m pan.t
```

Critical Section Example (cont.)

```
[root@moonzoo spin_test]# spin -t -p crit.pml
Starting P with pid 0
Starting P with pid 1
Starting Invariant with pid 2
      proc 1 (P) line 5 "crit.pml" (state 1)
                                                [(!(lock))]
  1:
  2:
      proc 0 (P) line 5 "crit.pml" (state 1) [(!(lock))]
      proc 1 (P) line 5 "crit.pml" (state 2) [lock = 1]
  3:
  4:
      proc 1 (P) line 6 "crit.pml" (state 3)
                                                [cnt = (cnt+1)]
       1 is in the crt sec!
  5: proc 1 (P) line 7 "crit.pml" (state 4)
                                                [printf('%d is in the crt sec!\\n',_pid)]
  6:
      proc 0 (P) line 5 "crit.pml" (state 2)
                                                [lock = 1]
  7: proc 0 (P) line 6 "crit.pml" (state 3) [cnt = (cnt+1)]
    0 is in the crt sec!
  8: proc 0 (P) line 7 "crit.pml" (state 4) [printf('%d is in the crt sec!\\n',_pid)]
spin: line 13 "crit.pml", Error: assertion violated
spin: text of failed assertion: assert((cnt<=1))
  9: proc 2 (Invariant) line 13 "crit.pml" (state 1) [assert((cnt<=1))]
spin: trail ends after 9 steps
#processes: 3
           lock = 1
           cnt = 2
  9:
      proc 2 (Invariant) line 14 "crit.pml" (state 2) <valid end state>
  9:
      proc 1 (P) line 8 "crit.pml" (state 5)
KOISTproc 0 (P) line 8 "crit.pml" (state 5)
3 processes created
```

```
bool lock;
                                      [root@moonzoo revised]# a.out
byte cnt;
                                      Full statespace search for:
                                          never claim
                                                              - (none specified)
active[2] proctype P() {
                                          assertion violations
                                                               +
     atomic{ !lock -> lock=true;}
                                          acceptance cycles - (not selected)
     cnt=cnt+1;
                                          invalid end states +
     printf("%d is in the crt sec!\n",_pid);
     cnt=cnt-1;
                                     State-vector 36 byte, depth reached 14, errors: 0
     lock=false;
                                         62 states, stored
                                         17 states, matched
                                         79 transitions (= stored+matched)
                                          0 atomic steps
active proctype Invariant() {
                                     hash conflicts: 0 (resolved)
     assert(cnt \leq 1);
                                     4.879
                                             memory usage (Mbyte)
```

Deadlocked Critical Section Example

```
[[root@moonzoo deadlocked]# a.out
                                       pan: invalid end state (at depth 3)
bool lock;
                                       (Spin Version 4.2.7 -- 23 June 2006)
byte cnt;
                                       Warning: Search not completed
                                            + Partial Order Reduction
active[2] proctype P() {
     atomic{ !lock -> lock==true;}
                                       Full statespace search for:
     cnt=cnt+1:
                                            never claim
                                                               - (none specified)
     printf("%d is in the crt sec!\n",_pid); assertion violations +
     cnt=cnt-1;
                                            acceptance cycles - (not selected)
     lock=false;
                                            invalid end states
                                                                  ÷
                                       State-vector 36 byte, depth reached 4, errors: 1
                                           5 states, stored
active proctype Invariant() {
                                           0 states, matched
     assert(cnt <= 1);
                                           5 transitions (= stored+matched)
                                           2 atomic steps
                                       hash conflicts: 0 (resolved)
```

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4.879 memory usage (Mbyte)

[root@moonzoo deadlocked]# spin -t -p deadlocked_crit.pml Starting P with pid 0

Starting P with pid 1

Starting Invariant with pid 2

1: proc 2 (Invariant) line 13 "deadlocked_crit.pml" (state 1) [assert((cnt<=1))]

2: proc 2 terminates

3: proc 1 (P) line 5 "deadlocked_crit.pml" (state 1) [(!(lock))]

4: proc 0 (P) line 5 "deadlocked_crit.pml" (state 1) [(!(lock))]

spin: trail ends after 4 steps

#processes: 2

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lock = 0cnt = 0

4: proc 1 (P) line 5 "deadlocked_crit.pml" (state 2)

```
4: proc 0 (P) line 5 "deadlocked_crit.pml" (state 2)
```

3 processes created

Communication Using Message Channels

- Spin provides communications through various types of message channels
 - Buffered or non-buffered (rendezvous comm.)
 - Various message types
 - Various message handling operators

Syntax

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- 4 chan ch1 = [2] of { bit, byte};
 - ch1!0,10;ch1!1,20
 - ch1?b,bt;ch1?1,bt Sender (1,20) (0,10) +Receiver

4 chan ch2= [0] of {bit, byte}

Basic channel inquiry

- 4 len(ch)
- \rm 4 empty(ch)
- \rm full(ch)
- \rm hempty(ch)
- 4 nfull(ch)

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Additional message passing operators

- 4 ch?[x,y]: polling only
- **ch**?<**x**,**y**>: copy a message without removing it
- 4 ch!!x,y: sorted sending (increasing order)
- 4 ch??5,y: random receiving
- 4 ch?x(y) == ch?x,y (for user's understandability)
- Be careful to use these operators inside of expressions
 - They have side-effects, which spin may not allow

Faulty Data Transfer Protocol

(pg 27, data switch model proposed at 1981 at Bell labs) mtype={ini,ack, dreq,data, shutup,quiet, dead}

chan M = [1] of {mtype}; chan W = [1] of {mtype};



The Sieve of Eratosthenes (pg 326)

```
proctype sieve(chan c; int prime)
/*
                                                        chan child = [0] of { mtype, int };
  The Sieve of Eratosthenes (c. 276-196 BC)
                                                        bool haschild; int n;
  Prints all prime numbers up to MAX
                                                        printf("MSC: %d is prime\n", prime);
*/
                                                   end: do
#define MAX
                25
                                                        :: c?number(n) ->
mtype = { number, eof };
                                                              if
chan root = [0] of { mtype, int };
                                                              :: (n%prime) == 0 -> printf("MSC: %d
                                                   = %d*%d\n", n, prime, n/prime)
init
                                                              :: else ->
     int n = 2;
                                                                   :: !haschild -> /* new prime */
     run sieve(root, n);
                                                                        haschild = true;
     do
                                                                         run sieve(child, n);
     :: (n < MAX) \rightarrow n++; root!number(n)
                                                                   :: else ->
     :: (n \ge MAX) \ge root!eof(0); break
                                                                        child!number(n)
     od
                                                                   fi;
                                                             fi
                                                        :: c?eof(0) -> break
                                                        od;
                                                        if
                                                        :: haschild -> child!eof(0)
                                                        :: else
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                                                        fi
```

Simulation Run

[moonzoo@verifier spin]\$ spin sieve-of-eratosthenes.pml

```
2 MSC: 2 is prime
   3 MSC: 3 is prime
  MSC: 4 = 2^{2}
     5 MSC: 5 is prime
  MSC: 6 = 2^{*}3
  MSC: 8 = 2*4
        7 MSC: 7 is prime
     MSC: 9 = 3^{*}3
  MSC: 10 = 2*5
  MSC: 1/2 = 2*6
  MSC: 14 = 2*7
          11 MSC: 11 is prime
     MSC: 15 = 3*5
             13MSC: 13 is prime
  MSC: 16 = 2*8
  MSC: 18 = 2*9
  MSC: 20 = 2*10
```