Propositional Calculus - *Semantics* (2/3)

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Overview

- 2.1 Boolean operators
- 2.2 Propositional formulas
- 2.3 Interpretations
- 2.4 Logical equivalence and substitution
- 2.5 Satisfiability, validity, and consequence
- 2.6 Semantic tableaux
- 2.7 Soundness and completeness

Logical equivalence

- Defn 2.13. Let $A_1, A_2 \in \mathcal{F}$. If $\nu(A_1) = \nu(A_2)$ for all interpretation, then A_1 is logically equivalent to A_2, denoted $A_1 \equiv A_2$
- Example 2.14. Is $p \vee q$ equivalent to $q \vee p$?

p	q	$\nu(p \vee q)$	$\nu(q \vee p)$
T	*T*	*T*	*T*
T	*F*	*T*	*T*
F	*T*	*T*	*T*
F	*F*	*F*	*F*

Logical equivalence

- We can extend the result of example 2.14 from atomic propositions to general formulas
- Theorem 2.15 Let A_1 and A_2 be any formulas. Then $A_1 \vee A_2 \equiv A_2 \vee A_1$.
  - Proof
    - Let ν be an arbitrary interpretation for $A_1 \vee A_2$. Then, ν is an interpretation for $A_2 \vee A_1$, too.
    - Similarly, ν is an interpretation for A_1 and A_2
    - Therefore, $\nu(A_1 \vee A_2)$=T iff $\nu(A_1)$ =T or $\nu(A_2)$ =T iff $\nu(A_2 \vee A_1)$=T

Logical equivalence

Definition 2.22

- A binary operator o is defined from a set of operators $\{o_1, \ldots o_n\}$ if and only if there is a logical equivalence $A_1 \circ A_2 \equiv A$, where A is a formula constructed from occurrences of A_1 and A_2 using the operator $\{o_1, \ldots, o_n\}$.
- Similarly, the unary operator $\neg$ is defined from a set of operators $\{o_1, \ldots o_n\}$ iff $\neg A_1 \equiv A$, where A is constructed from occurrences of A_1 and the operators in the set.
- Examples
  - $\leftrightarrow$ is defined from $\{\rightarrow, \wedge\}$ because $A \leftrightarrow B \equiv (A \rightarrow B) \wedge (B \rightarrow A)$
  - $\rightarrow$ is defined from $\{\neg, \vee\}$ because $A \rightarrow B \equiv \neg A \vee B$
  - $\wedge$ is defined from $\{\neg, \vee\}$ because $A \wedge B \equiv \neg(\neg A \vee \neg B)$

Object language v.s. metalanguage

- Note that '≡' is not a binary operator used in propositional logic (object language).
- '≡' (metalanguage) is used to explain a relationship between two formulas.
- Theorem 2.16
  - $A_1 \equiv A_2$ if and only if $A_1 \leftrightarrow A_2$ is true in every interpretation

Logical substitution

- Logical equivalence justifies substitution of one formula for another
- Defn 2.17 *A* is subformula of *B* if the formation tree for *A* occurs as a subtree of the formation tree for *B*. *A* is proper subformation of *B* if *A* is a subformation of *B*, but *A* is not identical to *B*.
- Example 2.18 The formula $(p \rightarrow q) \leftrightarrow (\neg p \rightarrow \neg q)$ contains the following proper subformulas:

  $p \rightarrow q$, $\neg p \rightarrow \neg q$, $\neg p$, $\neg q$, p and q

Logical substitution

- Def. 2.19
  - If A is a subformula of B and A' is any formula,
  - then B', the substitution of A' for A in B, denoted $B\{A \leftarrow A'\}$, is the formula obtained by replacing all occurrences of the subtree for A in B by the tree for A'.
- Theorem 2.21 Let A be a subformula of B and let A' be a formula such that $A \equiv A'$. Then $B \equiv B\{A \leftarrow A'\}$
- One of the most important applications of substitution is simplication
  - Ex. $p \land (\neg p \lor q) \equiv (p \land \neg p) \lor (p \land q) \equiv \text{false} \lor (p \land q) \equiv p \land q$

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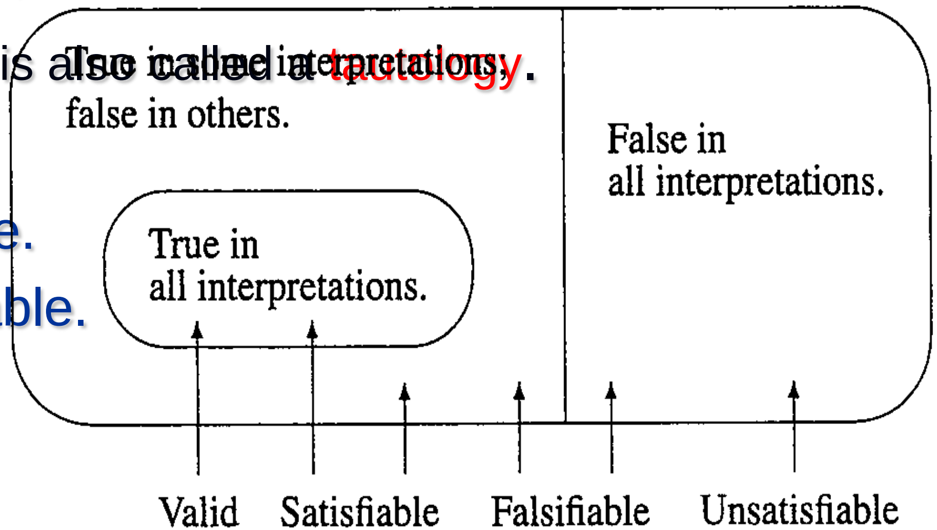
Satisfiability v.s. validity

- Definition 2.24
  - A propositional formula *A* is satisfiable iff $\nu(A)=T$ for some interpretation ν.
    - A satisfying interpretation is called a model for *A*.
  - *A* is valid, denoted $\models A$, iff $\nu(A) = T$ for all interpretation ν.
    - A valid propositional formula is also called a tautology.
- Theorem 2.25
  - *A* is valid iff $\neg A$ is unsatisfiable.
  - *A* is satisfiable iff $\neg A$ is falsifiable.



Satisfiability v.s. validity

Definition 2.26

- Let $\mathcal{V}$ be a set of formulas. An algorithm is a **decision procedure** for $\mathcal{V}$ if given an arbitrary formula A $\in \mathcal{F}$, it terminates and return the answer 'yes' if A $\in \mathcal{V}$ and the answer 'no' if A $\notin \mathcal{V}$
- By theorem 2.25, a decision procedure for satisfiability can be used as a decision procedure for validity.
  - Suppose $\mathcal{V}$ is a set of all satisfiable formulas
  - To decide if A is valid, apply the decision procedure for satisfiability to $\neg$A
    - This decision procedure is called a refutation procedure

Satisfiability v.s. validity

- Example 2.27 Is $(p \rightarrow q) \rightarrow (\neg q \rightarrow \neg p)$ valid?

p	q	$p \rightarrow q$	$\neg q \rightarrow \neg p$	$(p \rightarrow q) \rightarrow (\neg q \rightarrow \neg p)$
T	*T*	*T*	*T*	*T*
T	*F*	*F*	*F*	*T*
F	*T*	*T*	*T*	*T*
F	*F*	*T*	*T*	*T*

- Example 2.28 $p \vee q$ is satisfiable but not valid

Logical consequence

- Definition 2.30 (extension of satisfiability from a single formula to a set of formulas)
  - A set of formulas $U = \{A_1 , \ldots A_n\}$ is (simultaneously) satisfiable iff there exists an interpretation ν such that $\nu(A_1) = \ldots = \nu(A_n) = T$.
  - The satisfying interpretation is called a model of U.
  - U is unsatisfiable iff for every interpretation ν, there exists an i such that $\nu(A_i) = F$.

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Logical consequence

- Let U be a set of formulas and A a formula. If A is true in every model of U, then A is a logical consequence of U.
  - Notation: $U \models A$
  - If U is empty, logical consequence is the same as validity
- Theoem 2.38
  - $U \models A$ iff $\models A_1 \wedge \ldots \wedge A_n \rightarrow A$ where $U=\{A_1 \ldots A_n\}$
  - Note Theorem 2.16
    - $A_1 \equiv A_2$ if and only if $A_1 \leftrightarrow A_2$ is true in every interpretation

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Theories

- Logical consequence is the central concept in the foundations of mathematics
  - Valid formulas such as $p \vee q \leftrightarrow q \vee p$ are trivial and not interesting
  - Ex. Euclid assumed five formulas about geometry and deduced an extensive set of logical consequences.
- Definition 2.41
  - A set of formulas $\mathcal{T}$ is a theory if and only if it is **closed** under logical consequence.
    - $\mathcal{T}$ is closed under logical consequence if and only if for all formula A, if $\mathcal{T} \models A$ then $A \in \mathcal{T}$.
  - The elements of $\mathcal{T}$ are called theorems
- Let U be a set of formulas. $\mathcal{T}(U) = \{A \mid U \models A\}$ is called the theory of U. The formulas of U are called axioms and the theory $\mathcal{T}(U)$ is axiomatizable.
  - Is $\mathcal{T}(U)$ theory?

Examples of theory

- $U = \{ p \vee q \vee r, q \to r, r \to p \}$
- Interpretation v_1, v_3 and v_4 are models of U
- Which of the followings are true?
  - $U \vDash p$
  - $U \vDash q \to r$
  - $U \vDash r \vee \neg q$
  - $U \vDash p \wedge \neg q$
- Theory of U, i.e, $\mathcal{T}(U)$
  - $U \subseteq \mathcal{T}(U)$
    - because for all formula $A \in U$, $A \vDash A$
  - $p \in \mathcal{T}(U)$
    - because $U \vDash p$
  - $q \to r \in \mathcal{T}(U)$
    - because $U \vDash q \to r$
  - $p \wedge (q \to r) \in \mathcal{T}(U)$
    - because $U \vDash p \wedge (q \to r)$
      - since $U \vDash p$ and $U \vDash q \to r$ $\therefore$
  - …

	p	q	r	$p \vee q \vee r$	$q \to r$	$r \to p$
v_1	T	T	T	T	T	T
v_2	T	T	F	T	F	T
v_3	T	F	T	T	T	T
v_4	T	F	F	T	T	T
v_5	F	T	T	T	T	F
v_6	F	T	F	T	F	T
v_7	F	F	T	T	T	F
v_8	F	F	F	F	T	T

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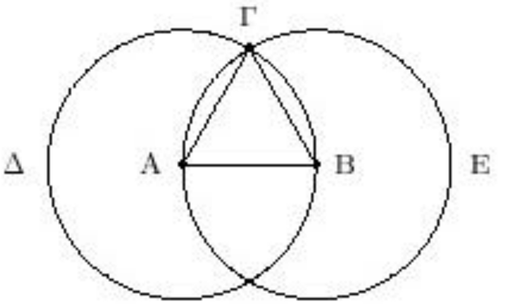
Ex. Theory of Euclidean geometry

- A set of 5 axioms U = {A_1,A_2,A_3,A_4,A_5} such that
  - A_1:Any two points can be joined by a unique straight line.
  - A_2:Any straight line segment can be extended indefinitely in a straight line.
  - A_3:Given any straight line segment, a circle can be drawn having the segment as radius and one endpoint as center.
  - A_4:All right angles are congruent.
  - A_5:For every line *l* and for every point P that does not lie on *l* there exists a unique line *m* through P that is parallel to *l*.
- Euclidean theory $\mathcal{T}_{Euclid} = \mathcal{T}(U) = \{ A \mid U \models A\}$
  - I.e.,$\mathcal{T}_{euclid}$ is axiomatizable by the above 5 axioms
  - Ex. one logical consequence of the axioms
    - given a line segment AB, an equilateral triangle exists that includes the segment as one of its sides.



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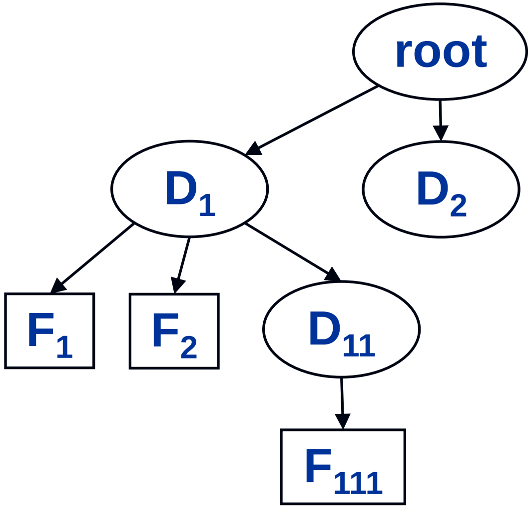
Ex2. Model checking (formal verification)

- A file system M can be specified by the following 7 formulas (i.e., a file system model M = { $A_1, A_2, A_3, A_4, A_5, A_6, A_7$ })
  - A_1:A file system object has one or no parent.
    - sig FSObject { parent: lone Dir }
  - A_2:A directory has a set of file system objects
    - sig Dir extends FSObject { contents: set FSObject }
  - A_3:A directory is the parent of its contents
    - fact defineContents { all d: Dir, o: d.contents | o.parent = d }
  - A_4: A file in the file system is a file system object
    - sig File extends FSObject {}
  - A_5: All file system objects are either files or directories
    - fact fileDirPartition { File + Dir = FSObject }
  - A_6: There exists only one root
    - one sig Root extends Dir { }{ no parent }
  - A_7: File system is connected
    - fact fileSystemConnected { FSObject in Root.*contents }



- We can prove that this file system does not have a cyclic path
  - A: No cyclic path exists
    - assert acyclic { no d: Dir | d in d.^contents }
  - $M \models A$ (i.e., this file system M does not have cyclic path)